



# 离散数学 (011122)



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- 2.1 Basic Concepts of Propositional Logic
- 2.2 Propositional logic equivalence transformations
- 2.3 Logical normal form

- **2.1.1 Propositions and Connectives**
  - Propositions and Truth Values (Simple Propositions, Compound Propositions)
  - Connectives ( $\neg$ ,  $\wedge$ ,  $\vee$ ,  $\rightarrow$ ,  $\leftrightarrow$ )
  
- **2.1.2 Propositional Formulas and Their Classification**
  - Propositional Formulas and Their Assignments
  - Truth Tables
  - Classification of Propositional Formulas

- **Proposition:** A statement that can be judged as true or false.
- **Truth value of a proposition:** The result of the judgment, either true or false.
- **True proposition:** A proposition with a truth value of true.
- **False proposition:** A proposition with a truth value of false.

 **Note:** Exclamatory sentences, imperative sentences, and interrogative sentences are not propositions. Also, paradoxes in declarative sentences or those with indeterminate judgment results are not propositions.

e.g. >>> Example: Which of the following sentences are propositions?

(1) The capital of the People's Republic of China is Beijing. True proposition

(2)  $2 + 3 = 6$ .

False proposition

(3)  $x + y > 8$ .

Undecided True Value

(4) Can you play tennis?

Interrogative sentences

(5) There is life on planets other than Earth.. Indeterminate judgment results

(6) This !

Exclamatory sentences

(7) Please close the door!

Imperative sentences

(8) I am lying.

Paradoxes

**(1), (2), (5) are propositions. (3), (4), (6)~(8) are not propositions.**

- **Simple Proposition (Atomic Proposition):** A proposition formed by simple statements.
  - **Symbolization of Simple Propositions:** Symbolization of simple propositions: Represented by  $p, q, r, \dots, p_i, q_i, r_i$  (where  $i \geq 1$ ). '1' represents true, and '0' represents false.
- **Compound Proposition:**  
A statement formed by connecting simple propositions with logical connectives.  
**Examples:**
  - e.g.* >>>
  - (1) If the weather is good tomorrow, we will play ball.  
Let  $p$ : The weather will be good tomorrow, and  $q$ : We will play ball. **If  $p$ , then  $q$ .**
  - (2) Luffy is drinking milk tea and scrolling through his phone.  
Let  $p$ : Luffy is drinking milk tea, and  $q$ : Luffy is scrolling through his phone.  **$p$  and  $q$ .**

### ↳ 2.1.1 Propositions and Connectives • $\neg$ and $\wedge$

- **Definition 2.1:** "Non- $p$ " (or "the negation of  $p$ ") is called the *negation* of  $p$ , denoted as  $\neg p$ . The symbol  $\neg$  is the negation connective, and it is defined such that  $\neg p$  is true if and only if  $p$  is false.

#### Example:

- $p$ : 2 is a composite number.  $\neg p$ : 2 is not a composite number.
- Since 2 is actually a prime number,  $p$  is false, and therefore  $\neg p$  is true.

- **Definition 2.2:**

" $p$  and  $q$ " (or " $p$  with  $q$ ") is called the *conjunction* of  $p$  and  $q$ , denoted as  $p \wedge q$ . The symbol  $\wedge$  is the conjunction connective, and it is defined such that  $p \wedge q$  is true if and only if both  $p$  and  $q$  are true simultaneously.

#### Example:

- $p$ : 2 is an even number.  $q$ : 2 is a prime number.
- $p \wedge q$ : 2 is an even prime number.
- Since 2 is indeed both even and prime, both  $p$  and  $q$  are true, so  $p \wedge q$  is also true

Symbolize the following propositions.

- |   |   |
|---|---|
| (1) Wang Xiao is smart and hardworking.                         | (1) $p \wedge q$<br>(Let $p$ : Wang Xiao is smart, $q$ : Wang Xiao is hardworking.)                           |
| (2) Wang Xiao is not only smart but also hardworking.           | (2) $p \wedge q$  |
| (3) Wang Xiao is smart, but not hardworking.                    | (3) $p \wedge \neg q$   |
| (4) Wang Xiao is not unintelligent, but rather not hardworking. | (4) $\neg (\neg p) \wedge \neg q$   |
| (5) Both Zhang Hui and Wang Li are outstanding students.        | (5) $r \wedge s$<br>(Let $r$ : Zhang Hui is an outstanding student, $s$ : Wang Li is an outstanding student.) |
| (6) Zhang Hui and Wang Li are classmates                        | (6) $t$ (A simple proposition, "and" connects two nouns, the entire sentence is a simple proposition.)        |

### ↳ 2.1.1 Propositions and Connectives • $\vee$

- **Definition 2.3:** " $p$  or  $q$ " is called the disjunctive form of  $p$  and  $q$ , denoted as  $p \vee q$ . The symbol  $\vee$  is called *the Disjunction connective*, and  $p \vee q$  is false if and only if both  $p$  and  $q$  are false.

*e.g.* >>> Example: Wang Yan has studied English or French.

Let  $p$ : Wang Yan has studied English,

$q$ : Wang Yan has studied French.

Symbolized as  $p \vee q$ .

### ↳ 2.1.1 Propositions and Connectives • Inclusive or vs. Exclusive

- The English term for " $\vee$ " is "inclusive" (Inclusive), corresponding to the "or" in everyday language. The English term for " $\wedge$ " is "exclusive" (Exclusive), corresponding to the "and" in everyday language. "Inclusive or" and "exclusive or" are their combined forms.

*e.g.* >>> Example: This task is to be done by either Zhang San or Li Si.  
Let  $p$ : Zhang San does this task,  $q$ : Li Si does this task. It should be symbolized as:

$$(1) (p \wedge \neg q) \vee (\neg p \wedge q).$$

$$(2) p \oplus q$$

$$(3) P \text{ XOR } q$$

e.g. >>> Example: Symbolize the following propositions:

- (1) 2 is a prime number or 4 is a prime number.
- (2) 2 is a prime number or 3 is a prime number.
- (3) 4 is a prime number or 6 is a prime number.

**Solve:** Let:

$p$ : 2 is a prime number,

$q$ : 3 is a prime number,

$r$ : 4 is a prime number,

$s$ : 6 is a prime number.

$$(1) p \vee r, \quad 1 \vee 0 = 1$$

$$(2) p \vee q, \quad 1 \vee 1 = 1$$

$$(3) r \vee s, \quad 0 \vee 0 = 0$$

### ↳ 2.1.1 Propositions and Connectives • $\vee$ (e.g.)

e.g. >>> Example: Symbolize the following propositions:

(4) Yuan Yuan can take an apple or a pear.

**Solve:** Let:

$t$ : Yuan Yuan takes an apple,

$u$ : Yuan Yuan takes a pear.

$$(t \wedge \neg u) \vee (\neg t \wedge u)$$

(5) Wang Xiaohong was born in 1975 or 1976

**Solve:** Let:

$v$ : Wang Xiaohong was born in 1975,

$w$ : Wang Xiaohong was born in 1976.

$$(v \wedge \neg w) \vee (\neg v \wedge w)$$

Can logical expression  
“ $(v \wedge \neg w) \vee (\neg v \wedge w)$ ”  
be converted to  
“ $v \vee w$ ”?

#### ■ Definition 2.4

"If  $p$ , then  $q$ " is called the implication of  $p$  and  $q$ , denoted as  $p \rightarrow q$ .  $p$  is called the antecedent (or hypothesis) of the implication, and  $q$  is called the consequent (or conclusion). The symbol  $\rightarrow$  is called the *implication connective*.

- It is defined that  $p \rightarrow q$  is false if and only if  $p$  is true and  $q$  is false.

#### e.g. >>> Example:

"If the weather is good tomorrow, we will go on an outing."

Let  $p$ : The weather is good tomorrow,

$q$ : We will go on an outing.

This can be formalized as  $p \rightarrow q$ .

### ↳ 2.1.1 Propositions and Connectives • → (Cont'd)

- **Logical relationship of  $p \rightarrow q$ :**  
 $q$  is a necessary condition for  $p$ , and  $p$  is a sufficient condition for  $q$ .
- **Various ways to express "If  $p$ , then  $q$ " (all having the same truth value as  $p \rightarrow q$ ):**
  - (1) If  $p$ , then  $q$ .
  - (2) If  $p$ , just  $q$ . (Emphasizes that the existence of  $p$  is a necessary condition for the existence of  $q$ .)
  - (3)  $p$  only if  $q$ . (Means that for  $p$  to be true,  $q$  must also be true.)
  - (4) Only if  $q$ , then  $p$ . (Indicates that the truth of  $q$  is a necessary condition for the truth of  $p$ .)
  - (5) Unless  $q$ , not  $p$ . (If  $q$  is not the case ( $q$  is false), then  $p$  does not hold ( $p$  is false).)

#### ■ Truth and Falsity of the Implication $p \rightarrow q$ :

- When  $p$  is true,  $q$  must also be true for  $p \rightarrow q$  to be "true".
- When  $p$  is true and  $q$  is false,  $p \rightarrow q$  is "false".
- When  $p$  is false, the implication does not impose any restrictions on the truth value of  $q$ , and  $p \rightarrow q$  is "true".

#### e.g. >>> Example:

Let  $p$  be defined as "It is raining today," and  $q$  be defined as "The ground is wet." The implication  $p \rightarrow q$  can be described as "If it rains today, then the ground will be wet."

- In this case, the implication is **only false** if it rains ( $p$  is true) and the ground is not wet ( $q$  is false). In all other scenarios (such as it not raining, or it raining and the ground being wet), the implication is true.

#### e.g. >>> Example:

Let  $p$ : It is cold, Let  $q$ : Wang wears a down jacket.,  
Here's the translation of the propositions with logical symbolism:

- (1) As long as it is cold, Xiao Wang will wear a down jacket.  $p \rightarrow q$
- (2) Because it is cold, Xiao Wang wears a down jacket.  $p \rightarrow q$
- (3) If Xiao Wang does not wear a down jacket, then it is not cold.  $\neg q \rightarrow \neg p$  or  $p \rightarrow q$
- (4) Only if it is cold, Xiao Wang will wear a down jacket.  $q \rightarrow p$

**e.g. >>> Example:**

Let  $p$ : It is cold, Let  $q$ : Wang wears a down jacket.

Here's the translation of the propositions with logical symbolism:

(5) Unless it is cold, Xiao Wang will wear a down jacket.

$$q \rightarrow p$$

(6) Unless Xiao Wang wears a down jacket, otherwise it will not be cold.

$$p \rightarrow q$$

(7) If it is not cold, then Xiao Wang will not wear a down jacket.

$$\neg p \rightarrow \neg q \text{ or } q \rightarrow p$$

(8) Xiao Wang wears a down jacket only if it is cold.

$$q \rightarrow p$$

#### ■ Definition 2.5:

The statement " $p$  if and only if  $q$ " is called the equivalence of  $p$  and  $q$ , denoted by  $p \leftrightarrow q$ , where  $\leftrightarrow$  is called the *biconditional operator*. It is defined that  $p \leftrightarrow q$  is true if and only if both  $p$  and  $q$  are true or both are false.

#### ■ Logical Relationship of $p \leftrightarrow q$ :

$p$  and  $q$  are mutually sufficient and necessary conditions for each other.

#### *e.g.* $\ggg$ Example:

Zhang San can do this task well, and only Zhang San can do it well.

Let  $p$ : Zhang San does the task,  $q$ : The task is done well.

This can be formalized as:  $p \leftrightarrow q$ .

e.g. **Example:** Determine the truth value of the following compound propositions:

- |   |   |
|---|---|
| (1) $2 + 2 = 4$ if and only if $3 + 3 = 6$ .  | 1 |
| (2) $2 + 2 = 4$ if and only if 3 is even.   | 0 |
| (3) $2 + 2 = 4$ if and only if the sun rises in the east.   | 1 |
| (4) $2 + 2 = 5$ if and only if the sun rises in the west.   | 0 |
| (5) A necessary and sufficient condition for $f(x)$ to be differentiable at $x_0$ is that it is continuous at $x_0$ . | 0 |

#### ■ Truth Values of Basic Compound Propositions

$p$	$q$	$p \wedge q$	$p \vee q$	$p \rightarrow q$	$p \leftrightarrow q$
0	0	0	0	1	1
0	1	0	1	1	0
1	0	0	1	0	0
1	1	1	1	1	1

#### ■ Precedence of Logical Connectives:

Parentheses ( $()$ ), Negation ( $\neg$ ), Conjunction ( $\wedge$ ), Disjunction ( $\vee$ ), Implication ( $\rightarrow$ ), Biconditional ( $\leftrightarrow$ ).

- **Same Level:** Evaluated from left to right.

- **Propositional Constants:** Simple propositions
- **Propositional Variables:** Variables that can take the value 0 (true) or 1 (false).
- **Definition 2.6 Well-Formed Formula (Propositional Formula, Formula):** A well-formed formula is **recursively defined** as follows:
  - (1) A single propositional constant or variable is a well-formed formula, also called an atomic formula.
  - (2) If A is a well-formed formula, then  $(\neg A)$  is also a well-formed formula.
  - (3) If A and B are well-formed formulas, then  $(A \wedge B)$ ,  $(A \vee B)$ ,  $(A \rightarrow B)$ , and  $(A \leftrightarrow B)$  are also well-formed formulas.
  - (4) Only those expressions formed by a finite number of applications of (1) through (3) are considered well-formed formulas.

e.g.  $\gggg$  Example:  $0$ ,  $p$ ,  $\neg p \vee q$ ,  $(p \vee q) \wedge (\neg p \vee r)$ ,  $p \vee q \rightarrow r$ ,  $(p \rightarrow q) \rightarrow r$

#### ■ Definition 2.7

(1) A single propositional variable or propositional constant is a 0-layer formula.

(2) A formula  $A$  is an  $(n+1)$ -layer formula (where  $n \geq 0$ ) if one of the following conditions is met:

- ①  $A = \neg B$ , where  $B$  is an  $n$ -layer formula.
- ②  $A = B \wedge C$ , where  $B$  and  $C$  are  $i$ -layer and  $j$ -layer formulas, respectively, and  $n = \max(i, j)$ .
- ③  $A = B \vee C$ , where  $B$  and  $C$  are  $i$ -layer and  $j$ -layer formulas, respectively, and  $n = \max(i, j)$ .
- ④  $A = B \rightarrow C$ , where  $B$  and  $C$  are  $i$ -layer and  $j$ -layer formulas, respectively, and  $n = \max(i, j)$ .
- ⑤  $A = B \leftrightarrow C$ , where  $B$  and  $C$  are  $i$ -layer and  $j$ -layer formulas, respectively, and  $n = \max(i, j)$ .

*e.g.* >>> **Example:** The propositional formulas

$p$  (0-layer)

$\neg p$  (1-layer)

$\neg p \rightarrow q$  (2-layer)

$(\neg(p \rightarrow q)) \leftrightarrow r$  (3-layer)

$((\neg p \wedge q) \rightarrow r) \leftrightarrow (\neg r \vee s)$  (4-layer)

#### ■ Definition 2.8

- Let  $p_1, p_2, \dots, p_n$  be all the propositional variables appearing in the formula  $A$ .
- Assigning a set of truth values to  $p_1, p_2, \dots, p_n$  is called an **assignment** or **interpretation** for  $A$ .
- An assignment that makes the formula true is called a ***satisfying assignment***, and an assignment that makes the formula false is called a ***falsifying assignment***.

#### ■ Definition 2.8 Explanation:

(1) An assignment is denoted as  $a = a_1 a_2 \dots a_n$ , where each  $a_i$  is either 0 or 1, and the  $a_i$  are written without any punctuation marks between them.

(2) Generally, the assignment corresponds to the propositional variables in the order of their subscripts or alphabetical order. That is:

- When all propositional variables in  $A$  are  $p_1, p_2, \dots, p_n$ , assigning  $a_1 a_2 \dots a_n$  to  $A$  means  $p_1 = a_1, p_2 = a_2, \dots, p_n = a_n$ .
- When all propositional variables in  $A$  are  $p, q, r, \dots$ , assigning  $a_1 a_2 a_3 \dots$  to  $A$  means  $p_1 = a_1, p_2 = a_2, p_3 = a_3, \dots$ .

e.g. >>> **Example:**

■ Formula  $A = (\neg p_1 \wedge \neg p_2 \wedge \neg p_3) \vee (p_1 \wedge p_2)$

- The assignment 000 is a **satisfying assignment** (makes the formula **true**).
- The assignment 001 is a **falsifying assignment** (makes the formula **false**).

■ Formula  $B = (p \rightarrow q) \rightarrow r$

- The assignment 000 is a **falsifying assignment** (makes the formula **false**).
- The assignment 001 is a **satisfying assignment** (makes the formula **true**).

- **Truth Table:** A list of the values taken by a propositional formula under all possible assignments.
    - A formula with  $n$  variables has  $2^n$  assignments.
- e.g. >>> **Example:** Provide the truth table for the following propositional formula.

(1)  $(q \rightarrow p) \wedge q \rightarrow p$

$p$	$q$	$q \rightarrow p$	$(q \rightarrow p) \wedge q$	$(q \rightarrow p) \wedge q \rightarrow p$
0	0	1	0	1
0	1	0	0	1
1	0	1	0	1
1	1	1	1	1

e.g. >>> **Example:** Provide the truth table for the following propositional formula.

$$(2) \neg (\neg p \vee q) \wedge q$$

$p$	$q$	$\neg p$	$\neg p \vee q$	$\neg (\neg p \vee q)$	$\neg (\neg p \vee q) \wedge q$
0	0	1	1	0	0
0	1	1	1	0	0
1	0	0	0	1	0
1	1	0	1	0	0

e.g. >>> Example: Provide the truth table for the following propositional formula.

(3)  $(p \vee q) \rightarrow \neg r$

$p$	$q$	$r$	$p \vee q$	$\neg r$	$(p \vee q) \rightarrow \neg r$
0	0	0	0	1	1
0	0	1	0	0	1
0	1	0	1	1	1
0	1	1	1	0	0
1	0	0	1	1	1
1	0	1	1	0	0
1	1	0	1	1	1
1	1	1	1	0	0

- **Tautology (Always True):** A propositional formula that is never false under any assignment.
- **Contradiction (Always False):** A propositional formula that is never true under any assignment.
- **Satisfiable Formula:** A propositional formula that is not a contradiction.

**Note:** A tautology is satisfiable, but the converse is not true.

#### Examples:

- e.g.* >>> (1)  $(q \rightarrow p) \wedge q \rightarrow p$  is a tautology.
- (2)  $\neg(\neg p \vee q) \wedge q$  is a contradiction.
- (3)  $(p \vee q) \rightarrow \neg r$  is a satisfiable formula that is not a tautology.